Normal edge-colorings of cubic graphs

Vahan Mkrtchyan October 23, 2023

Outline

- 1. Some Conjectures in Graph Theory
- 2. Normal edge-colorings of cubic graphs
- 3. Our result
- 4. Future work

Collaborators

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Some Conjectures in Graph Theory

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- The elements of E are called edges of (the graph) G.

An example of a graph

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If $V = \{a, b, c, 1, 2\}$ and $E = <\{1, 2\}, \{a, b\}, \{1, 2\}, \{b, c\}, \{a, c\}, \{b, c\} >$, then G = (V, E) is a graph.

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Simple and parallel edges of G

 $\{1,2\}$ and $\{b,c\}$ have multiplicity two in *G*. They are called multi-edges or parallel edges of *G*. Other edges have multiplicity one and they are called simple edges of *G*.

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- The degree of a vertex is the number of edges incident to it.



Figure 2: Degrees of vertices of G = (V, E).

Some more definitions

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Figure 3: An example of a graph G = (V, E) with $\chi'(G) = 4$.

Maximum Degree and Multiplicity of a Graph

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For any graph G: $\Delta(G) \leq \chi'(G) \leq \Delta(G) + \mu(G)$.

Corollary (Vizing)

For any simple graph G (i.e. $\mu(G) = 1$): $\Delta(G) \le \chi'(G) \le \Delta(G) + 1$.



Figure 4: An example of a graph G = (V, E) with $\Delta(G) = 3$, $\mu(G) = 2$ and $\chi'(G) = 4$.



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Corollary

If G is a cubic graph, then

 $3 \leq \chi'(G) \leq 4.$



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Theorem (Holyer)

The problem of deciding whether a given cubic graph G has $\chi'(G) = 3$ is NP-complete.

Independent edges and Matchings

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Figure 5: An example of a matching.

An example of a perfect matching in a graph

An example

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Figure 6: A perfect matching in K₄.

Bridges in graphs

An edge e of a graph G is a bridge, if G - e has more connected components than G does.

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Figure 7: *e* is a bridge in a cubic graph *G*.

Cubic graphs and perfect matchings

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Observation

If G is a cubic graph, then

$$\chi'(G) = 3$$
 if and only if $k(G) = 3$.

The Petersen graph

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Figure 8: The graph P_{10} .

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The Petersen graph is the smallest bridgeless cubic graph G with $\chi'(G) = 4$.

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- $k(P_{10}) = 5$.

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Conjecture (Berge-Fulkerson)

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For any bridgeless cubic graph G we have

$$\chi'(2G)=6.$$

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Theorem (G. Mazzuoccolo, 2011)

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Theorem (G. Mazzuoccolo, 2011)

Conjectures of Berge and Berge-Fulkerson are equivalent.

Definitions

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- For i = 1, ..., t let $l(C_i)$ be the number of edges of C_i , and let $l(\mathcal{C}) = \sum_{i=1}^{t} l(C_i)$.

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- For i = 1, ..., t let $l(C_i)$ be the number of edges of C_i , and let $l(\mathcal{C}) = \sum_{i=1}^{t} l(C_i)$.
- $I(\mathscr{C})$ is called the length of the cycle cover \mathscr{C} .

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Conjecture ((5,2) **Even Subgraph Cover Conjecture)**

Any bridgeless (not necessarily cubic) graph G has a cycle cover $\mathscr{C} = (C_1, ..., C_t)$, such that each edge of G belongs to exactly 2 of the cycles of \mathscr{C} .

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Any bridgeless (not necessarily cubic) graph G has a cycle cover $\mathscr{C} = (C_1, ..., C_t)$, such that $l(\mathscr{C}) \leq \frac{7}{5} \cdot |E|$.

(5,2) Even Subgraph Cover Conjecture implies Cycle Double Cover Conjecture.

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Shortest Cycle Cover Conjecture implies Cycle Double Cover Conjecture.

The relationship
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Figure 9: The relationship among the five conjectures.

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Definition

If G and H are two cubic graphs, then an H-coloring of G is a mapping $f : E(G) \to E(H)$ such that for each vertex $x \in V(G)$, there is a vertex $y \in V(H)$ with $f(\partial_G(x)) = \partial_H(y)$.

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Definition

If G admits an H-coloring f, then we will write $H \prec G$.

An example of an *H*-coloring

An example: $H \prec G$

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Figure 10: An example of an *H*-coloring of *G*.

If $H \prec G$ and $K \prec H$, then $K \prec G$.

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3-edge-colorable cubic graphs

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If $H \prec G$ and f is an H-coloring of G, then for any adjacent edges e and e' of G, one has: f(e) and f(e') are adjacent in H.

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Adjacent edges: Opposite statement

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Adjacent edges: Opposite statement

If *H* is triangle-free, and a mapping $f : E(G) \to E(H)$ has the property that for any adjacent edges *e* and *e'* of *G*, f(e) and f(e') are adjacent in *H*,

If $H \prec G$ and $K \prec H$, then $K \prec G$.

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If *H* is triangle-free, and a mapping $f : E(G) \to E(H)$ has the property that for any adjacent edges *e* and *e'* of *G*, f(e) and f(e') are adjacent in *H*, then *f* is an *H*-coloring of *G*.

Suppose that G and H are cubic graphs with $H \prec G$, and let f be an H-coloring of G. Then:

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- (a) If M is any matching of H, then $f^{-1}(M)$ is a matching of G;
- (b) $\chi'(G) \leq \chi'(H)$, where $\chi'(G)$ is the chromatic index of G;
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- (d) $k(G) \leq k(H)$.

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- (d) $k(G) \leq k(H)$.
- (e) For every even subgraph C of H, $f^{-1}(C)$ is an even subgraph of G;
- (f) For every bridge e of G, the edge f(e) is a bridge of H.

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- (d) $k(G) \leq k(H)$.
- (e) For every even subgraph C of H, $f^{-1}(C)$ is an even subgraph of G;
- (f) For every bridge e of G, the edge f(e) is a bridge of H.
- (g) If H is bridgeless, then G is also bridgeless.

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Conjecture (P₁₀-conjecture, 1988)

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Conjecture (*P*₁₀-conjecture, 1988)

For any bridgeless cubic graph G, one has: $P_{10} \prec G$.

*P*₁₀-conjecture implies Berge-Fulkerson conjecture.

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Observation

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P₁₀-conjecture implies Shortest Cycle Cover Conjecture.

The relationship

The relationship among the six conjectures

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Figure 11: The relationship among the six conjectures.

Normal edge-colorings of cubic graphs

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If a cubic graph G admits a P_{10} -coloring f, such that for an edge e of P_{10} $f^{-1}(e) = \emptyset$, then G is 3-edge-colorable.

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A snark is a bridgeless cubic graph that is not 3-edge-colorable.

Observation

If a cubic graph G admits a P_{10} -coloring f, such that for an edge e of P_{10} $f^{-1}(e) = \emptyset$, then G is 3-edge-colorable.

Corollary

If G is a snark, then any P_{10} -coloring of G must use all edges of P_{10} .

Poor, rich edges and normal edge-colorings of cubic graphs

Let f be a k-edge-coloring of a cubic graph G, and let $S_f(w)$ be the set of colors of edges of G incident to the vertex w.

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A *k*-edge-coloring *f* of a cubic graph *G* is normal, if any edge of *G* is poor or rich in *f*. Let $\chi_N(G)$ be the smallest *k* for which *G* admits a normal *k*-edge-coloring. Clearly, $\chi_N(G) \ge 3$ for any cubic graph *G*. $\chi_N(G) = 3$, if and only if $\chi'(G) = 3$. An example

An example of a normal edge-coloring of a cubic graph

An example



Figure 12: A cubic graph that requires 7 colors in a normal coloring.
An example of a normal edge-coloring of a cubic graph

An example



Figure 12: A cubic graph that requires 7 colors in a normal coloring.

The bridge is poor.

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The bridge is poor. All other edges are rich.

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Figure 12: A cubic graph that requires 7 colors in a normal coloring.

The bridge is poor. All other edges are rich. It can be shown that $\chi_N(G) = 7$.

Does any cubic graph admit a normal k-edge-coloring for some k?

Does any cubic graph admit a normal k-edge-coloring for some k?

An example

Does any cubic graph admit a normal k-edge-coloring for some k?

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Figure 13: An example of a cubic graph that does not admit a normal coloring.

A cubic graph G admits a P_{10} -coloring (i.e. $P_{10} \prec G$), if and only if $\chi'_N(G) \leq 5$.

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Conjecture (*P*₁₀**-conjecture restated)**

For any bridgeless cubic graph G, we have $\chi'_N(G) \leq 5$.

Our result

Question

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Summary of prior results

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Theorem (Our main result)

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Theorem (Our main result)

For any simple cubic graph G, we have $\chi'_N(G) \leq 7$.

Future work

Improving our result

 Proving χ'_N(G) ≤ 7 in the class of bridgeless cubic graphs is relatively easy (8-flow theorem).

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Conjecture (Intermediate conjecture)

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Conjecture (Intermediate conjecture)

For any bridgeless cubic graph G, we have $\chi'_N(G) \leq 6$.

THANK YOU!